Challenges and Solutions for Fast Remote Persistent Memory Access

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We finally have fast durable storage

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How should we build distributed systems for NVMM?

Recent DRAM-based systems for fast networks provide a blueprint

- Key-value stores: Pilaf [ATC 13], MICA [NSDI 2014], HERD [SIGCOMM 2014], ...
- Transaction processing systems: FaRM [NSDI 14, SOSP 15], DrTM [SOSP 15, OSDI 18], FaSST [OSDI 16], NAM-DB [VLDB 17], ...
- State machine replication: DARE [HPDC 2015], Zookeeper-in-a-box [NSDI 2016], ...

What design decisions need to change if we use NVMM instead of DRAM?
The power-safe domain in NVMM systems

![Diagram of NVMM system](image)

- Core 1
- Core N
- L3 cache
- clwb
- NVMM
- CPU
Latency of Remote Persistent Memory Access

Two approaches:

- Remote Procedure Calls (RPCs)
- Remote Direct Memory Access (RDMA)

Finding: RDMA has similar as RPCs for durable writes to remote NVMM
NVMM writes with Remote Procedure Calls (RPCs)

![Diagram showing the flow of data between Client and Server through NIC, L3, and NVMM with RPC request and response]]
NVMM writes with Remote Direct Memory Access (RDMA)
NVMM with RDMA, Direct Cache Access (DCA) disabled
NVMM with Remote Direct Memory Access (RDMA)
NVMM removes latency advantage of RDMA over RPCs

**Critical path of persistent RPC:**
One network RTT + one PCIe RTT

**Critical path of persistent RDMA:**
One network RTT + one PCIe RTT
NVMM removes latency advantage of RDMA over RPCs

Experiment setup:
- Cascade Lake Xeon CPUs
- 6x 256 GB Optane DIMMs
- 56 Gbps ConnectX-3 InfiniBand
Application: State Machine Replication

Network messages in Raft SMR

Client latency (median, µs)

- Client: 6.6µs
- Leader: 8.1µs
- RDMA (best case)

RPCs

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Bandwidth of large writes to remote NVMM

Client goodput (Gbps)

- RDMA (DCA): 45 Gbps
- RPCs: 22 Gbps

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RDMA: Direct Cache Access randomizes NVMM write order

Root: CPU cache line size (64 bytes) < Optane memory erase block size (256 bytes)
Solution: Disable Direct Cache Access

- Several systems enable Direct Cache Access when using NVMM with RDMA
- Disabling a hardware optimization (Direct Cache Access) improves efficiency!
Bandwidth of large writes to remote NVMM

![Graph showing bandwidth comparison between RDMA (DCA) and RPCs]

- RDMA (DCA): 45 Gbps
- RPCs: 22 Gbps
Problem: CPU cores are slow at writing to NVMM

DMA engine improves bulk write bandwidth with RPCs from 22 Gbps to 48 Gbps
Conclusions

• Designing fast networked systems for NVMM requires attention to new low-level factors

• Our techniques can help while the hardware improves:
  • Better mechanisms for durable RDMA
  • Direct Cache Access without NVMM access order randomization
  • Faster persistent copying for CPU cores
  • See paper for more!

Thank you!